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$$
\begin{equation*}
\frac{2}{\frac{1}{a}+\frac{1}{b}}<\sqrt{a b}<\frac{b-a}{\ln b-\ln a}<\frac{a+b}{2}<\sqrt{\frac{a^{2}+b^{2}}{2}} \quad(0<a<b) \tag{4}
\end{equation*}
$$

relating the harmonic, geometric, logarithmic, arithmetic, and root-mean-square means.

For a rigorous proof of (1), observe that

$$
\cosh x-1 \equiv\left(\frac{e^{x / 2}-e^{-x / 2}}{2}\right)^{2}>0 \quad \text { for } \quad x>0
$$

Therefore,

$$
\frac{1}{\cosh ^{2} x}<1<\cosh x \quad \text { for } \quad x>0
$$

and (1) follows by integrating the preceding inequality from $x=0$ to $x=t$. The other inequalities of (3) are obvious.

## Trees and Tennis Rankings

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Everyone is familiar with the usual kind of "elimination" tennis tournament in which losers drop out and winners continue to play until one undefeated player remains. Such a tournament is nicely described by a tree whose vertices $1,2, \ldots, n$ represent the $n$ players and whose edges represent the matches that were playedthe direction of each edge pointing toward the loser. In the tournament of Figure 1, for example, player 4 won by beating player 6 (who beat 11 and 1), player 10, and player 3 (who beat 8,2 , and 5 ). In addition, player 8 beat player 9 (who beat 7 and 12).


Figure 1.
Suppose we wanted to rank the players of this tournament, subject to the rule: for each edge of $T$, player i precedes player $j$ if and only if $i$ beats $j$. There is clearly no unique way to do so since $T$ does not contain enough information for a full account
of the players' relative strengths. Indeed, although 4 (the root of tree $T$ ) must begin the list, the next member of the ranking could be 6,10 , or 3 . Therefore, the general question to be asked is:

How many rankings can be generated by a given rooted (directed) tree?
Let us first show that for the tournament in Figure 1, the answer is 158400 . The analysis of the general case is simple and pretty.


$T_{c}$
Figure 2a.


Figure 2b.

Let (Figure 2a) $a, b, c$ be the players who lost to 4 . Then $a, b, c$ are the roots of subtrees $T_{a}, T_{b}, T_{c}$ of $T$. Knowing the number of rankings $R\left(T_{a}\right), R\left(T_{b}\right), R\left(T_{c}\right)$ for these subtrees will enable us to compute the number of rankings $R(T)$ for $T$ itself.

There $\operatorname{are}\left(\begin{array}{rrr} & 11 & \\ 3 & 1 & 7\end{array}\right)=\frac{11}{3!1!7!}$ ways to string 11 beads using 3 amber, 1 blue, and 7 crimson beads. Assume that these colored beads are precisely the numbered vertices of $T_{a}, T_{b}, T_{c}$. Then for any particular color arrangement on the string, there are respective $R\left(T_{a}\right), R\left(T_{b}\right), R\left(T_{c}\right)$ rankings for each color. (Figure 2 b depicts one possible color arrangement and one possible ranking for that given color arrangement.) Thus,

$$
R(T)=\frac{11!}{3!1!7!} \times R\left(T_{a}\right) \times R\left(T_{b}\right) \times R\left(T_{c}\right)
$$

Clearly, $R\left(T_{a}\right)=2$ and $R\left(T_{b}\right)=1$. To obtain $R\left(T_{c}\right)$, we can invoke the preceding argument, applied to subtrees $T_{8}, T_{2}, T_{5}$ of $T_{c}$. This yields

$$
R\left(T_{c}\right)=\frac{6!}{4!1!1!} \times 2 \times 1 \times 1 .
$$

All together, we obtain

$$
R(T)=\frac{11!}{3!1!7!} \times 2 \times 1 \times\left(\frac{6!}{4!1!1!} \times 2 \times 1 \times 1\right)=158400
$$

Turning now to the general case, we assume that the tournament-representing tree $T$ can be arranged so that the root $r$ occurs by itself at the top, in "level 0 ," of the tree (Figure 3). Let the players $w_{1}, w_{2}, \ldots, w_{k}$ who lost to winner $r$ be placed (in any order) in level 1 of the tree.


Figure 3.
Each vertex $w_{i}$ is the root of a subtree $T_{w_{i}}$ of $T$; each $T_{w_{i}}$ represents a subtournament, the ranking of whose players provides us with a problem that is identical to the kind of problem we started with. Let us show that, if we can solve the problem for the subtrees $T_{w_{i}}$, we can succeed with $T$ itself.

Every ranking for $T$ has a place for each of the $n$ players and, because the first position must go to winner $r$, our freedom extends only through the remaining $n-1$ positions. We don't care how the vertices of the various subtrees get interlaced, so long as the vertices of each $T_{w_{i}}$, considered alone, satisfy " $i$ before $j$ means $i$ beat $j$." Consequently, a ranking for $T$ is constructed by the double procedure:
(1) For each $i$, select $\left|T_{w_{i}}\right|$ places among the $n-1$ open positions for the vertices of $T_{w_{i}}$;
(2) Enter in these places the $\left|T_{w_{i}}\right|$ vertices of $T_{w_{i}}$ in any acceptable $R\left(T_{w_{i}}\right)$ ranked order.

Practically by definition, the number of ways of achieving (1) is the multinomial coefficient

$$
\binom{n-1}{\left|T_{w_{1}}\right|\left|T_{w_{2}}\right| \cdots\left|T_{w_{k}}\right|}=\frac{(n-1)!}{\left|T_{w_{1}}\right|!\left|T_{w_{2}}\right|!\cdots\left|T_{w_{k}}\right|!} .
$$

And for each of these partitions, the number of ways of doing (2) is $R\left(T_{w_{1}}\right) \times$ $R\left(T_{w_{2}}\right) \times \cdots \times R\left(T_{w_{k}}\right)$. Therefore, the total number of rankings for $T$ is given by the recurrence relation

$$
R(T)=\frac{(n-1)!}{\left|T_{w_{1}}\right|!\left|T_{w_{2}}\right|!\cdots\left|T_{w_{k}}\right|!} \cdot R\left(T_{w_{1}}\right) R\left(T_{w_{2}}\right) \cdots R\left(T_{w_{k}}\right)
$$

with the obvious initial value $R(T)=1$ for $|T|=1$.

